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SIGNAL PROCESSING APPARATUS, PROGRAM FOR USE IN SIGNAL  
PROCESSING APPARATUS, STORAGE MEDIUM STORING THEREON PROGRAM  
FOR USE IN SIGNAL PROCESSING APPARATUS, AND SIGNAL  
PROCESSING METHOD

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates to a signal processing apparatus, a program for use in a signal processing apparatus, a storage medium storing thereon a program for use in a signal processing apparatus, and a signal processing method; for converting a continuous signal such as image information or audio information into other information by means of digital conversion such as filtering or level adjusting.

2. Description of the Related Art

In a conventional signal processing apparatus designed to be used for such a purpose, when data obtained as a result of an operation performed on given input data has a longer bit length than that of the input data, it is necessary to adjust the resultant data so as to have the same bit length as that of the input data.

Fig. 6 is a diagram illustrating an example of a process for the above-described purpose. In Fig. 6, P<sub>n</sub> is input data that has been converted into digital form, and Q<sub>n</sub>

is data according to which to control  $P_n$ . In general,  $Q_n$  is less than 1.  $P_n$  and  $Q_n$  are input to a multiplier 1, and  $X_n (= P_n \times Q_n)$  is output. In general,  $X_n$  has a bit length equal to the sum of the bit length of input data  $P_n$  and the bit length of  $Q_n$ .

To avoid a problem with the difference in bit length between the output data  $X_n$  and the input data  $P_n$ , it is necessary to convert the output data  $X_n$  so as to have the same bit length as that of the input data  $P_n$ .

To this end,  $X_n$  is input to a thresholder 2 to shorten the bit length. As a result, output data  $Y_n$  having a properly adjusted bit length is obtained. Note that the bit length of the output data  $Y_n$  does not necessarily need to be the same as that of the input data  $P_n$ .

The thresholding described above can be performed by (1) unconditional rounding down, (2) unconditional round up, or (3) rounding off (rounding up or down depending on a given value).

For example, when  $X$  is given as an input value, if operation  $Y = 0.02X + 1$  is performed on  $X$  and the resultant value of  $Y$  is rounded down or rounded off to an integer, the output value is obtained as shown in Fig. 7 (the output value obtained by means of rounding up is not shown in Fig. 7).

In the case of rounding off, the integral value

obtained by rounding off  $Y$  is 1 for the input value smaller than 25, but the integral value jumps to 2 when the input value  $X$  becomes 25. As can be seen from Fig. 7, the integral value obtained by rounding off or rounding down the calculated value  $Y$  changes in a stepwise fashion (in a discontinuous fashion).

In the above-described conventional technology, the final output value includes a more or less error whichever technique of (1) rounding down, (2) rounding up, or (3) rounding off is used. In the case of rounding off, if errors are averaged over a long period, the averaged error becomes 0. In contrast, when rounding-down is used, each output value includes an error of 0.5 on average. Even in the case of rounding off, if output values are not averaged, a great error occurs for values close to each discontinuous transition point. In the case in which the input value  $X$  varies linearly as in the example described above, the error alternately increases and decreases as the input value  $X$  varies. However, when the input value  $X$  varies within a narrow range close to a transition point, or when the input value  $X$  varies at a low rate, the rounded value of  $Y$  includes a large error and remains unchanged for a long period.

Furthermore, when there is a range difference in bit length between the input data and the output data, the

cumulative value of the output data can have a great cumulative error with respect to the cumulative value of the input data. When a large cumulative error causes a problem, signal processing is performed using a large number of bits such that the error arising via the signal processing is minimized and such that the cumulative error does not cause the output data to be biased in a particular direction.

Whether the error causes a problem or not depends on the bit length of the output data and on for what purpose the output data is used. In general, the bit length used in the signal processing is determined such that the error falls within an allowable range.

That is, the error can be reduced to a sufficiently small range by increasing the bit length used in the signal processing or by increasing the memory capacity.

However, in some cases, it is not allowed to increase the bit length to reduce the error (that is, in some applications, it is required to perform signal processing using a limited number of bits). For example, in some applications, only on-off processing is allowed. In some other applications, on-off processing is more desirable to achieve a high energy efficiency although continuous processing is also allowed.

In some applications such as level adjustment of digital audio data or contrast adjustment of digital image

data, if the dynamic range is reduced as a result of level adjustment or contrast adjustment, the reduction in dynamic range causes a reduction in the number of bits that can be used to represent signals in the reduced dynamic range becomes smaller. This means that the relative error with respect to the signal increases as the signal level decreases.

In thermal ink-jet printers, a fixed amount of ink droplet is emitted at a time and the amount of ink droplet emitted at a time cannot be changed. That is, when an ink image is formed on a recording medium, emission of ink at each dot is performed by means of two-level control in which a fixed amount of ink droplet is emitted or no ink is emitted. To achieve halftone representation using two-level ink emission at each dot, the density of inked dots per unit area is modulated. To this end, error diffusion or dithering is widely used. However, to achieve good enough halftone representation, complicated computation is needed, and thus a complicated circuit and a long computation time are required.

#### SUMMARY OF THE INVENTION

An object of the present invention is to provide a technique to perform signal processing without creating significant degradation in smoothness or frequency

characteristics, even if the bit length is reduced when an input digital signal is processed.

In an aspect, to achieve the above object, the present invention provides a signal processing apparatus for receiving digital signals that are continuously related and input sequentially, performing a predetermined operation on each of sequentially input digital signals, and outputting a result of the operation, the signal processing apparatus comprising operation means for performing the predetermined operation on an input digital signal, high-order part extraction means for extracting a necessary high-order part by rounding off a result of the operation performed by the operation means, difference calculation means for calculating the difference between the result of the operation performed by the operation means and the high-order part extracted by the high-order part extraction means, and feedback means for adding, to a next input digital signal, the difference value calculated by the difference calculation means or a value obtained by performing a predetermined operation on the difference value calculated by the difference calculation means.

In this signal processing apparatus according to the present invention, an error arising via rounding-off processing is calculated by the difference calculation means, and the error or a value obtained by performing the

predetermined operation on the error is added to a next input digital signal via the feedback means.

Because each error is fed back in such a manner that an error arising via an operation performed on an input digital signal is added to a next input digital signal, an error arising via an operation performed on the next input digital signal is added to a further next input digital signal, and so on, accumulation of errors is prevented.

#### BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a block diagram illustrating a signal processing apparatus according to a first embodiment of the present invention;

Fig. 2 is a diagram illustrating, in the form of a graph, an operation result;

Fig. 3 is a graph showing results of operations (operation A and operation B) that are similar to the operation shown in Fig. 2 and that are performed simultaneously such that there is a phase difference between them;

Fig. 4 is a diagram showing examples of images formed by controlling the density of dots per unit area using the process according to the present invention;

Fig. 5 is a block diagram of a signal processing apparatus according to a second embodiment of the present

invention;

Fig. 6 is a diagram showing a data processing method according to a conventional technique; and

Fig. 7 is a graph showing the relationship between input and output values for a case in which the fractional portions of input values are discarded and for a case in which input values are rounded off to nearest integers.

#### DESCRIPTION OF THE PREFERRED EMBODIMENTS

The present invention is described in further detail below with reference to specific embodiments in conjunction with the accompanying drawings.

##### First Embodiment

Fig. 1 is a block diagram showing a signal processing apparatus according to a first embodiment of the present invention. Although in the present embodiment, the signal processing apparatus 10 shown in Fig. 1 is assumed to be realized by hardware (for example, in the form of a signal processing circuit), the signal processing apparatus 10 may also be realized by software. All registers such as a register 11 shown in Fig. 1 operate in synchronization with the same clock signal.

In the present embodiment, by way of example, each input data X is subjected to an operation (performed by



calculation means according to the present invention)  
according to the following equation:

$$Y = 0.02X + 1.$$

A high-order part of a value obtained as a result of the above operation is extracted by rounding off the value to an integer (by high-order part extraction means according to the present invention), and a residual fraction is dealt with as an error.

In Fig. 1, an N-bit digital signal is input to the input register 11. The digital signal input to this input register 11 is transferred to an input adder 12. To the input adder 12, an error generated in a previous operation is also input. The input digital signal and the input error are added by the input adder 12.

The addition result is output to a high-order bit register 13 and a low-order bit register 14.

As shown in Fig. 1, if K bits of the input N-bit digital signal  $X_n$  is output to the high-order bit register 13, then the number of bits output to the low-order bit register 14 is  $(N - K)$  bits.

Thereafter, it is determined whether the value output from the low-order bit register 14 is equal to or greater than 0.5. If the value output from the low-order bit register 14 is equal to or greater than 0.5, the value is rounded up to the high-order bit value by adding 1 to the

high-order bit adder 15. As a result, 1 is added to the value output from the high-order bit register 13.

When 1 is input to the high-order bit adder 15, the error resulting from the rounding up is input to the low-order bit adder 16 (by difference calculation means according to the present invention). For example, if the value output from the low-order bit register 14 is 0.7, then it is determined that the value output from the low-order bit register 14 is greater than 0.5, and 1 is input to the high order bit adder 15. In this specific case, the error resulting from the rounding up is -0.3, and -0.3 is input to the low-order bit adder 16.

When rounding up is performed, if the value output from the low-order bit register 14 is P, then

$$(1 - P) \times (-1)$$

is calculated and the resultant value is input to the low-order bit adder 16.

On the other hand, when it is determined that the value output from the low-order bit register 14 is smaller than 0.5, 1 is not input to the high order bit adder 15, but the value (fraction part) output from the low-order bit register 14 is directly input to the low-order bit adder 16 (by the difference calculation means according to the present invention).

Thus, rounding off to the extracted integer part is

performed via the above-described process.

A K-bit value output from the high order bit adder 15 is sent to an output register 17 and output as data  $Y_n$  therefrom.

On the other hand, the output value from the low-order bit adder 16 is input to the fraction processing register 18, which in turn converts the received fractional value into a value with a predetermined number of bits (a predetermined number of decimal places). The resultant value is input to the input adder 12 thereby feeding back the error (by feedback means according to the present invention).

The process described above is described in further detail below.

Herein, it is assumed that input values are given in the form of a sequence of continuous real numbers  $\{X_1, X_2, X_3, \dots\}$ .

Each of the real numbers  $X_1, X_2, \dots$  consists of an integer part A and a fraction part B. That is, each real number can be expressed as

$$X_n = A_n + B_n$$

Let  $Y_n$  be an output value obtained by rounding off  $X_n$  to an integer. For example, when  $X_1$  is given as

$$X_1 = A_1 + B_1$$

then  $Y_1$  is given by

$$Y_1 = [X_1]$$

$$= \begin{cases} A1 & (\text{when } B1 < 0.5) \\ A1+1 & (\text{when } 0.5 \leq B1 < 1) \end{cases}$$

where  $[Xn]$  denotes the operation of rounding off  $Xn$  to a particular digit (in the present embodiment, a fractional part is rounded off to an integer).

If  $Y1 = A1$  (that is, if rounding up does not occur), then

$$\begin{aligned} Y2 &= [X2 (= A2 + B2) + B1] \\ &= \begin{cases} A2 & (\text{when } B2 + B1 < 0.5) \\ A2+1 & (\text{when } 0.5 \leq B2 + B1 < 1) \end{cases} \end{aligned}$$

On the other hand, when  $Y1 = A1 + 1$  (that is, when rounding up occurs),

$$\begin{aligned} Y2 &= [X2 (= A2 + B2) + (B1 - 1)] \\ &= \begin{cases} A2 & (\text{when } B2 + B1 - 1 < 0.5) \\ A2+1 & (\text{when } 0.5 \leq B2 + B1 - 1 < 1) \end{cases} \end{aligned}$$

All  $Yn$  are determined one by one in the above-described manner.

In the present embodiment, an error created as a result of rounding off is handled (added) when a next value is input. For example, if rounding up occurs in the calculation of  $Y1$ ,  $Y1$  is given by

$$Y1 = A1 + 1$$

Herein, although  $Y1$  is greater than the actual input value  $X1$ , the error is subtracted when next value  $Y2$  is determined. Therefore, this error is completely cancelled

out in  $Y1 + Y2$ .

Thus, when  $Y1, Y2, \dots, Yn$  are determined one by one, the sum of thereof ( $Y1 + Y2 + \dots + Yn$ ) is given by  $SUM(Yn) = (A1 + 1^*) + (A2 + 1^*) + (A3 + 1^*) + \dots + (An + 1^*) + (B1 - 1^*) + (B2 - 1^*) + (B3 - 1^*) + \dots + (Bn - 1^*)$   
 $= X1 + X2 + X3 + \dots + Xn$

where  $1^*$  takes 1 only when a carry appears as a result of rounding off.

In the above equation,  $(A1 + 1^*)$  to  $(An + 1^*)$  are integer parts, and  $(B1 - 1^*)$  to  $(Bn - 1^*)$  are fraction parts.

In the above equation, if  $1^*$  takes 1 (that is, if a carry appears), for example, in term  $A1$ , then  $1^*$  takes 1 also in term  $B1$ . Therefore, when  $Y1, Y2, \dots, Yn$  are added together, errors arising from rounding off are all cancelled out (that is, errors are not accumulated).

Thus, the sum of  $Yn$  (output values) becomes equal to the sum of  $Xn$  (input values).

Thus, when data to be processed is continuous, theoretically, no error occurs if the processing is performed using a large enough number of bits.

Fig. 2 shows, in the form of a graph, the result obtained when

$$Y = 0.02X + 1$$

is calculated in the above-described manner. In Fig. 5, for the purpose of comparison, the result obtained when the

calculation is performed according to the conventional rounding-off scheme is also shown.

As shown in Fig. 2, in the simple rounding off scheme, the output does not change unless the fractional part of an input value reaches a threshold value. In contrast, in the scheme according to the present embodiment, each fractional value arising as a result of rounding off is added to a next input value, and thus the output value changes frequently. The frequency of the change is proportional to the difference between the input value and the output value, that is, the absolute value of the error. The output value, which changes back and forth between two values, can be smoothed by means of averaging using a filter or the like.

Although the error arising as a result of rounding off in each calculation can be twice the error that occurs in the conventional rounding off scheme, the moving average taken over a plurality of inputs  $X$  has a small error because the output changes when the input changes a in a small range. In Fig. 2, the moving average taken over successive five points is shown. This curve indicates that the moving average changes within a small range.

In the scheme according to the present embodiment, as described above, in a range in which an error becomes great, the output value frequently changes back and forth between two values, and thus an averaged value becomes equal to the

median of the two values. Furthermore, because the majority of error components lie in a high frequency range, it is possible to easily remove the high-frequency components (errors) from the output simply by passing the output through a simple (low-order) lowpass filter.

Another example of the process of feeding back an error from the fraction processing register 18 to the input adder 12 is described below.

In the case in which continuous digital signals are sequentially input, errors can be well handled by feeding back the error via the input adder 12 in the above-described manner.

However, when inputs are given in the form of  $X_n = A_n + B_n$ , if inputting is stopped, a new fraction part  $B_n$  does not occur, and a current fraction part  $B_n$  to be added to a next input value remains without being added until the next input value  $X_n$  is actually given.

However, in this case, the fraction part  $B_n$  has no correlation with the next input value, and thus preserving of the fraction part  $B_n$  is not only meaningless but can be harmful depending on the next input value  $X_n$ .

For example, when a next input value  $X_n$  is very close to  $Y_n$  and thus no error should appear in  $Y_n$ , if an error greater than 0.5 remains, the output value becomes different from the expected value (that is,  $X_n + 1$  or  $X_n - 1$  is

output).

One technique to avoid the above problem is to reset the error (to be added to a next input value) stored in the fraction processing register 18 to zero, when the next input value  $X_n$  becomes zero after the operation of a sequence of digital signals (a first set of digital signals) is completed.

In this technique, when a next sequence of digital signals (a second set of digital signals)  $X_n$  is input, the operation is started from an initialized state.

Contrary to the above, a value stored in the fraction processing register 18 may be used to create a particular effect on the next operation. For example, when two operations are simultaneously performed under the control of the same clock signal, results can be output such that they have a particular relationship (for example, results are output not at the same time but at periodically shifted times or positions).

Fig. 3 is a graph showing results of operations (operation A and operation B) that are similar to the operation shown in Fig. 2 and that are performed simultaneously such that there is a phase difference between them. In Fig. 3, a solid line represents an operation A, and a broken line represents an operation B ( $A + 0.5$ ).

Fig. 4 shows examples of images formed by controlling



the density of dots per unit area using the process described above. In Fig. 4, values varying in units of steps from 1 to 100 are input, and the input values are binarized by thresholding them with respect to 100. Each time the sum of fed-back errors becomes greater than 100, a dot is formed. In Fig. 4, a time axis is taken in a direction denoted by an arrow (from up to down). That is, dots are formed from up to down in Fig. 4. 50 sequences of dots are obtained via independent processing, and results are shown in Fig. 4 such that sequences of dots are shifted sequence by sequence in a horizontal direction.

In 50 sequences of dots shown on the left-hand side of Fig. 4, random values from 1 to 100 are given as 50 initial values (located in a horizontal line at the top).

On the other hand, in 50 sequences of dots shown on the right-hand side of Fig. 4, a repetition of 6 values (properly ordered integral multiples 16, that is, 16, 32, 48, 64, 80, and 96) is given as initial values.

As can be seen from Fig. 4, when a plurality of operations are performed in parallel as in image processing, it is possible to change the initial phases by controlling the initial values thereby achieving desirable effects.

## Second Embodiment

Fig. 5 is a block diagram showing a signal processing

apparatus (signal processing apparatus 10A) according to a second embodiment of the present invention.

In the second embodiment, unlike the first embodiment described above with reference to Fig. 1 in which each fractional value (error) is directly added to a next input value  $X_n$ , the fractional value is added via a multiplier 21. That is, the fractional value output from the low-order bit adder 16 is input to the multiplier 21 and multiplied by a constant  $\alpha$  (for example, 0.75) smaller than 1.

If  $X_n$  is not input for a long period, the fractional value is reduced at each operation interval, and thus the fractional value decreases toward zero. Therefore, when  $X_n$  is input after a long period during which no  $X_n$  is input, an error (a fraction) added to the input value  $X_n$  has become very small, and thus the operation is performed in a state substantially identical to the initial state, although the very small error added to the input value  $X_n$ .

In the present embodiment, because errors that occur during the operation are fed back such that each error is added to a next input value  $X_n$ , error components occurring during the operation lie in a high frequency range. If the output including the error components is oversampled at a frequency higher than twice the Nyquist frequency, the error components are shifted to a further higher frequency range. This makes it possible to easily remove the high frequency

components from the output by passing the output through a low-order lowpass filter without causing a significant loss of a desired signal component.

The signal processing apparatus 10 or 10A described above may be used not only to threshold a given signal with respect to a single reference value without creating a significant error, but also to threshold a given signal with respect to two or more reference values. Specific examples to which the signal processing apparatus according to the present invention can be applied include (1) power control (illumination control, temperature control, motion control, etc.), (2) measurement and indication of measurement result (meters), and (3) droplet emission control (ink-jet printers or the like).

The signal processing apparatus according to the present invention is essentially different from that based on the error diffusion method, as described below.

The error diffusion method is used mainly in image processing in which two-dimensional data is dealt with. In contrast, the signal processing method according to the present invention is basically used to process a sequence of one-dimensional data.

In the error diffusion method, errors are weighted and weighted errors are added to a plurality of adjacent data (that is, errors are diffused two-dimensionally).

In contrast, in the present invention, an error is simply added to a next input value. Furthermore, in the present invention, an initial value may be arbitrarily set, for example, by preloading the initial value.

Furthermore, unlike the error diffusion method, the present invention does not need a high-capacity memory, and operation can be performed in a short time.

Although the present invention has been described above with reference to specific embodiments, the present invention is not limited to the details of those embodiments, but various modifications are possible.

For example, rounding off can be performed at an arbitrary digit. Which part is to be extracted as a high-order part by means of rounding off may be set arbitrarily. Both the high-order part and the low-order part may be decimals. Conversely, both may be integers.

Although in the signal processing apparatuses and methods according to specific embodiments described above with reference to Fig. 1 or 5, it is assumed that those apparatuses or methods are implemented by means of hardware, apparatuses or methods according to the present invention may also be implemented by means of software. That is, the apparatus or the method described above with reference to Fig. 1 or 5 (note that the present invention is not limited to those shown in Figs. 1 and 5) may be achieved by

providing a program in an apparatus thereby implementing one or all of functions of the signal processing apparatus according to the present invention. The program for achieving one or all of functions of the signal processing apparatus according to the present invention may be stored on a storage medium such as a CD-ROM, and the CD-ROM may be provided to realize one or all of the functions of the signal processing apparatus according to the present invention.

As described above, the present invention makes it possible to improve the calculation accuracy in signal processing and reduce the calculation time, without needing either an increase in memory capacity or an increase in cost.